

The Specker-Blatter Theorem

In memory of Ernst P. Specker at the occasion of his centenary

J.A. Makowsky

Based on a survey paper by E. Fischer, T. Kotek and J.A. Makowsky
"Application of logic to combinatorial sequences and their recurrence relations."
Model Theoretic Methods in Finite Combinatorics 558 (2011): 1-42.

A video-taped longer version of this talk was given at the
Seminars "Proof Theory" and "Logic Online Seminar"
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http://www.mathnet.ru/php/seminars.phtml?option_lang=eng&presentid=27017

Abstract:

The Specker-Blatter Theorem states that number of labeled graphs over n vertices satisfying an MSOL-definable property satisfies a linear modular recurrence relation for every modulus m .

We give a gentle survey of the statement, the proof and applications.

[File:sp-intro](#)



Ernst Specker (1920-2011)



Christian Blatter (1935-2021)

C.B and E.S: Le nombre de structures finies d'une théorie à caractère fini,
Sciences Mathématiques, Fonds Nationale de la recherche Scientifique, Bruxelles, 1981

C.B and E.S: Modular periodicity of combinatorial sequences,
Abstract Amer. Math. Soc, 1983

C.B. and E.S.: Recurrence relations for the number of labeled structures on a finite set.
In Logic and Machines: Decision Problems and Complexity (pp. 43-61).
Springer Berlin/Heidelberg, 1984.

E.S.: Application of logic and combinatorics to enumeration problems
Trends in Theoretical Computer Science (E. Börger, ed.), 1988

scholar.google.com

- Ernst Speckers most cited paper, according to google scholar is:

Kochen, Simon and Specker, Ernst P.,
The problem of hidden variables in quantum mechanics,
reprinted in:

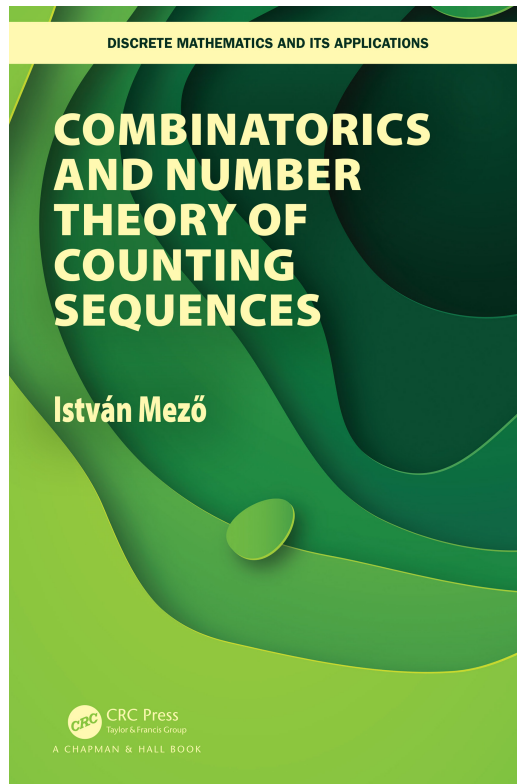
The logico-algebraic approach to quantum mechanics,
293–328, 1975 Springer

with 3163 citations

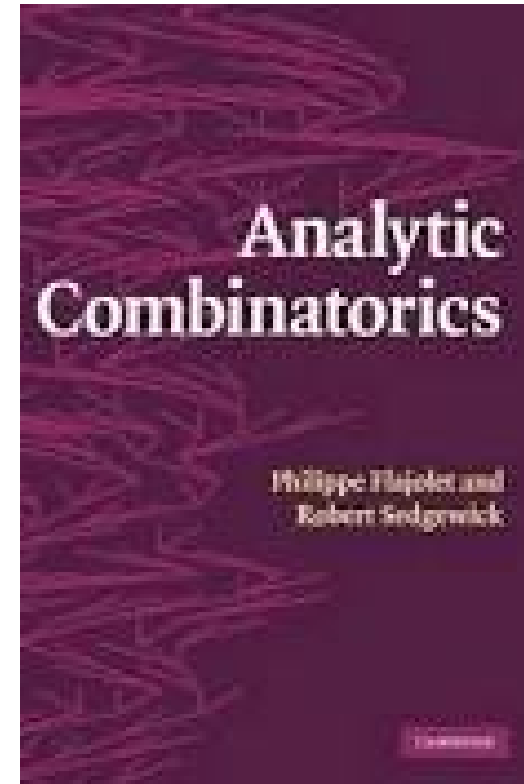
The talks after mine are about E. Specker's work in Quantum Mechanics

- I am going to talk about the [least cited paper](#).

Before I started working on it,
[all its version had together had less than 10 citations](#).



2020



2009

Even these books do not mention this work

Congruences for combinatorial functions 1980



Ira Gessel



Philippe Flajolet (1948-2011)

- Gessel, Ira.
Congruences for Bell and tangent numbers.
IBM Thomas J. Watson Research Division, 1978.
- Flajolet, Philippe. "On congruences and continued fractions for some classical combinatorial quantities."
Discrete Mathematics 41.2 (1982): 145-153.

Congruences for infinitely many combinatorial functions

- Let P be a graph property, and let

$$d_P(n) = |\{E \subseteq [n]^2 : ([n], E) \in P\}|$$

be the number of ways one can find an edge relation E on the set $\{1, \dots, n\}$ to obtain a graph $G = ([n], E) \in P$.

- $d_P(n)$ is called the **density function of P** .
- If P is monotone (hereditary) it is called the **speed of P** .

C. Blatter and E. Specker study
algebraic properties of density functions already in **1980**.

Only in **1994**, E. Scheinerman and J. Zito initiate the
study growth rates of density functions for hereditary properties,
and relate them to structural properties of graphs $G \in P$.

Comes the Hungarian School....

The Scheinerman-Zito paper started a new line of research:

- Scheinerman, Edward R., and Jennifer Zito. "On the size of hereditary classes of graphs." *Journal of Combinatorial Theory, Series B* 61.1 (1994): 16-39.
- Steger, Angelika. Asymptotic properties of H-free graphs. Diss. Habilitationsschrift, Forschungsinstitut für Diskrete Mathematik, Universität Bonn, 1993.
- Bollobás, Béla, and Andrew Thomason. "Hereditary and monotone properties of graphs." *The Mathematics of Paul Erdős II*. Springer, Berlin, Heidelberg, 1997. 70-78.
- Balogh, József, Béla Bollobás, and David Weinreich. "The speed of hereditary properties of graphs." *Journal of Combinatorial Theory, Series B* 79.2 (2000): 131-156.
- Balogh, József, Béla Bollobás, and Miklós Simonovits. "The number of graphs without forbidden subgraphs." *Journal of Combinatorial Theory, Series B* 91.1 (2004): 1-24.
- Balogh, József, Béla Bollobás, Michael Saks, and Vera T. Sós. "The unlabelled speed of a hereditary graph property." *Journal of Combinatorial Theory, Series B* 99, no. 1 (2009): 9-19.
- Alon, N., Balogh, J., Bollobás, B., and Morris, R. (2011). The structure of almost all graphs in a hereditary property. *Journal of Combinatorial Theory, Series B*, 101(2), 85-110.

The original theorem

(C. Blatter and E. Specker 1980)

Let P be a property of structures $(A, R_i), i \in [k]$, where

- (i) each R_i is a unary or binary relation,
- (ii) P is definable using Monadic Second Order Logic **MSOL**,

Let d_P the density function for P and d_P^m be the sequence

$$d_P^m(i) = d_P(i) \mod m.$$

Then $d_P^m(i) : i \in \mathbb{N}$ is ultimately periodic.

It was left open, whether

- (i) the restriction in (i) to at most binary relations was necessary, and
- (ii) the restriction in (ii) to **MSOL** among fragments of Second Order Logic **SOL** was necessary.

The density function for regular graphs

J.H. Redfield's Theorem, 1927

The class REG_r of simple regular graphs where every vertex has degree r is *FOL*-definable (for fixed r).

Counting the number of labeled regular graphs is treated completely in Chapter 7 of the book:

Harary, Frank, and Edgar M. Palmer. Graphical enumeration, 1973, reprinted by Elsevier, 2014.

where an explicit formula is given, essentially due to J.H. Redfield (1927) and rediscovered by R.C. Read (1959).

However, the formula is very complicated.

For cubic graphs, the function is explicitly given: $f_{\mathcal{R}_3}(2n+1) = 0$ and

$$f_{\mathcal{R}_3}(2n) = \frac{(2n)!}{6^n} \sum_{j,k} \frac{(-1)^j (6k-2j)! 6^j}{(3k-j)!(2k-j)!(n-k)!} 48^k \sum_i \frac{(-1)^i j!}{(j-2i)! i!}$$

Can you see that this is ultimately periodic mod 17?

Regular graphs of even degree

- Eulerian graphs are connected graphs of even degree.
- Graphs of even degree are not MSOL-definable, neither are Eulerian graphs.

Let CMSOL be the logic obtained from MSOL by adding modular counting quantifiers:

$$C_{m,a}(x)\phi(x)$$

which say that the set elements satisfying $\phi(x)$ in a finite structure $\mathfrak{A}$ has size a modulo m .

Regular graphs of even degree, and Eulerian graphs are CMSOL-definable.

What can we say about $d_{EULER}(n)$?

$d_{EULER}(n)$ is not a special case of the Blatter-Specker Theorem!

Two equal-sized cliques

Let EQ_2 the class of finite graphs which consist of the disjoint unions of two equal-sized cliques.

EQ_2 is not CMSOL-definable, but it is SOL-definable.

We look at its density function $d_{EQ_2}(n)$. We have

$$d_{EQ_2}(n) = \begin{cases} \frac{1}{2} \binom{2m}{m} & \text{for } n = 2m \\ 0 & \text{else} \end{cases}$$

The factor $\frac{1}{2}$ is there because we cannot distinguish the choice of the first clique from the choice of its complement.

Proposition: (Lucas, 1878)

For every n which is not a power of 2, we have $d_{EQ_2}(n) \equiv 0 \pmod{2}$, and for every n which is a power of 2 we have $d_{EQ_2}(n) \equiv 1 \pmod{2}$.

In particular, $d_{EQ_2}(n)$ is not ultimately periodic modulo 2.

A proof may be found as Exercise 5.61 in:
R. Graham, D. Knuth and O. Patashnik,
Concrete Mathematics, 2nd ed., Addison-Wesley 1994

E. Fischer's counter example

- From 2000 I started to lecture about the Specker-Blatter theorem.
- In 2002 Eldar Fischer, attending my lecture, found a counter example:

Theorem: There is a property of finite structures with one quaternary relation, which is even **FOL**-definable but violates the Blatter Specker Theorem.

Fischer, Eldar. "The SpeckerBlatter theorem does not hold for quaternary relations."
Journal of Combinatorial Theory, Series A 103.1 (2003): 121-136.

Recent work

E. Specker was intrigued and excited by Fischer's example.

- Specker, E.
"Modular Counting and Substitution of Structures."
Combinatorics, Probability & Computing 14.1-2 (2005): 203.

E. Fischer, and later my PhD student T. Kotek and I
analyzed the proof of the Blatter-Specker Theorem:

- Fischer, Eldar, and Johann A. Makowsky.
"The Specker-Blatter theorem revisited."
International Computing and Combinatorics Conference. Springer, Berlin,
Heidelberg, 2003.
- Kotek, Tomer, and Johann A. Makowsky. "Definability of combinatorial functions and their linear recurrence relations." Fields of logic and computation. Springer, Berlin, Heidelberg, 2010. 444-462.

And a more recent paper

Kotek, Tomer, and Johann A. Makowsky.

"On the Tutte and matching polynomials for complete graphs."

arXiv preprint arXiv:2112.06581 (2021).

submitted to the special issue of Fundamenta Informaticae
celebrating the centenary of Boris (Boaz) A. Trakhtenbrot

The Tutte polynomial

- The **Tutte polynomial** $T(G; X, Y)$ is defined as

$$T(G; X, Y) = \sum_{S \subseteq E(G)} (X - 1)^{\kappa(E) - \kappa(S)} \cdot (Y - 1)^{|S| - \kappa(S)}$$

where $\kappa(S)$ is the number of connected components of the spanning subgraph $G[S] = (V, S)$.

- I. Gessel (1979), and independently I. Pak (1993), proved:

$$T(K_n; a, b) = \sum_{k=1}^n \binom{n-1}{k-1} \left(a + \sum_{i=1}^{k-1} b^i \right) T(K_{k-1}; 1, b) \cdot T(K_{n-k}; a, b)$$

What can we say about $T(K_n; a, b)$?

What can we say about $T(K_n; a, b)$ modulo μ ?

Evaluations of $T(K_n; a, b)$

- $T(K_n; 2, 1)$ counts the number of forests on n vertices.
- $T(K_n; 1, 1)$ counts the number of trees on n vertices.
- $T(K_n; 1, 2)$ counts the number of connected graphs on n vertices.
- $C_n(b) = T(K_n; 1, b)$ has been analyzed by
A. P. Mani and R.J. Stones (2016) mod p^k .

A.P. Mani and R.J. Stones,
The number of labeled connected graphs modulo prime powers,
SIAM. J. Discrete Math., 30.2 (2016), pp. 1046-1057

They then formulate a conjecture for $T(K_n; a, b)$.

The conjecture by A. P. Mani and R.J. Stones

SIAM J. Discrete Math., vol. 30.2 (2016), 1046-1057

Let $\phi(m)$ be the [Euler totient function](#), which is defined as the number of integers $a \in [m]$ such that $\gcd(a, m) = 1$.

Let p be an odd prime and let k be a positive integer.

- If $n \geq p^k$, $a, b \in \mathbb{Z}$, and $b \not\equiv 1 \pmod{p}$, then

$$T(K_n; a, b) \equiv \begin{cases} b^{\frac{\phi(p)}{2}-1} & \text{if } p \geq 3, n = p \text{ and } a \equiv 1 \pmod{p} \\ b^{\frac{\phi(p^k)}{2}} \cdot T(K_{n-\phi(p^k)}; a, b) & \text{otherwise} \end{cases} \pmod{p}$$

- If $n \geq p^k$, $a, b \in \mathbb{Z}$, and $b \equiv 1 \pmod{p}$, then

$$T(K_n; a, b) \equiv \begin{cases} (n + a - 1)^{p^k} \cdot T(K_{n-p^k}; a, b) & n > p^k \\ (a - 1)^{p^k-1} & n = p^k \end{cases} \pmod{p^k}$$

Although, Mani and Stones cite our chapter on “Application of Logic to Combinatorial Sequences”, they misquote and obviously did not read it properly, missing out on the Specker-Blatter Theorem.

Inspired by the conjectures of
Mani and Stones we show

Theorem: (T. Kotek and JAM)

For every $a, b, \mu \in \mathbb{N}^+$ with $a, b > 1$,
 $\gcd(a - 1, \mu) = 1$ and $\gcd(b - 1, \mu) = 1$,
the sequence $T(K_n, a, b)$ is ultimately periodic modulo μ .

Theorem: (T. Kotek and JAM)

If P is a graph polynomial hypergraph-definable in
extended $\text{MSOL}_{\text{hyperg}}$ without order,
then for every $\mu \in \mathbb{N}$ and every $a, b \in \mathbb{Z}$
the sequence $P(K_n, a, b)$ is ultimately periodic modulo μ .

Variations on Stirling numbers of the second kind, I

Let $a, m \in \mathbb{N}^+$ and $A \subseteq \mathbb{N}^+$.

- $S(n, a)$ counts the number of partitions of $[n]$ into a non-empty sets.
- $S(n, a)$ is ultimately periodic modulo m since (for fixed a) the generating function is rational:

$$\sum_{n=0}^{\infty} S(n, a) x^n = \frac{x^a}{(1-x)(1-2x) \cdots (1-ax)}.$$

- Let $S_A(n, a)$ count the number of partitions of $[n]$ into a non-empty sets of cardinalities in A .
- $S_A(n, 1) = \begin{cases} 1 & n \in A \\ 0 & n \notin A \end{cases}$. Hence A is ultimately periodic iff $S_A(n, 1)$ is.

Variations on Stirling numbers of the second kind, II

Theorem: (JAM and V. Rakita)

For every fixed a, m the sequence $S_A(n, a)$ is ultimately periodic modulo m iff A is ultimately periodic.

I am confident, but not sure that there is an elementary direct proof of this with explicit description of the ultimate periodicity.

However, using the Specker-Blatter Theorem gives it **without computations**.

Computing $S(n, 2)$ and $S(n, 3)$

$S(n, 2)$ can be computed by counting all pairs of non-empty subsets $A, B \subset [n]$ of size k_A, k_B each. We note that $k_B = n - k_A$ and $k_A = n - k_B$ are both counted. We conclude that

$$S(n, 2) = 2^n - 2/2$$

But we can also count as follows

$$S(n, 2) = \frac{1}{2} \sum_{j=1}^{n-1} \binom{n}{j} = \sum_{n > j \geq n-j > 0} \binom{n}{j}$$

Similarly we get

$$S(n, 3) = \sum_{n > k_1 \geq k_2 \geq n - (k_1 + k_2) > 0} \binom{n}{k_1} \binom{n - k_1}{k_2}$$

Using the usual recurrence relation for $S(n, k)$ we get

$$S(n, 3) = \frac{1}{2}(3^{n-1} - 2^n + 1)$$

Checking for $n = 4$ we get with both methods $S(4, 3) = 6$.

Computing $S_A(n, a)$

We put $\ell_0 = 0$, $\ell_1 = k_1$ and $\ell_{i+1} = \ell_i + k_{\ell+1}$ and compute:

$$S_A(n, a) = \sum_{k_1 \geq \dots \geq k_a} \prod_{i=1}^a \binom{n - \ell_{i-1}}{k_i}$$

where $\sum_{i=1}^a k_i = n$ and $k_1, k_2, \dots, k_a \in A$.

Using the fact that

$$\binom{n}{a} \binom{n-k}{j} = \binom{n}{j} \binom{n-j}{k}$$

we see that the product in the above equation is independent of the order of the numbers $k_i \in A$. As we assume $k_1 \geq \dots \geq k_a \in A$ we sum over all ways n can be written as a sum of a summands in A **only once**.

Simplifying the product we get that for every $m \in \mathbb{N}^+$

$$S_A(n, a) = n! \cdot \sum_{k_1 \geq \dots \geq k_a \in A} \prod_{i=1}^a \frac{1}{k_i} = n! \cdot f(n)$$

for $f(n)$ not an integer, hence we cannot conclude that $n! \cdot f(n) \not\equiv 0 \pmod{m}$.

Example: $2^n = \sum_{j=0}^n \binom{n}{j} = \sum_{j=0}^n \frac{n!}{j! \cdot (n-j)!} = n! \cdot \sum_{j=0}^n \frac{1}{j! \cdot (n-j)!} \not\equiv 0 \pmod{3}$

My collaborators



Eldar Fischer



Tomer Kotek

- Fischer, Eldar, Tomer Kotek, and Johann A. Makowsky.
"Application of logic to combinatorial sequences and their recurrence relations.",
in: Model Theoretic Methods in Finite Combinatorics,
volume 558 of Contemporary Mathematics, (2011): 1-42.

Separating Logic from Combinatorics

C. Blatter and E. Specker associate with a property P of relational structures a rank, the **substitution rank**, with values in $\mathbb{N} \cup \{\infty\}$.

The logic part of their theorem is the following:

Proposition: If P is MSOL-definable this rank is finite.

We (EF and JAM) prove:

Proposition: If P is CMSOL-definable this rank is finite.

The rest of their (CB and ES) proof is algebra-combinatorial.

Proposition: If P is a property of structures with all relations at most binary, and P has finite rank, then $d_P^m(n)$ is ultimately periodic for every $m \in \mathbb{N}$.

This now covers also the case of **Eulerian graphs** and **regular graphs of even degree**.

Refining the analysis of the CB-ES proof

- We also take into account the

maximal degree of the structures in P .

- We generalize the Substitution rank of a property P using

Hankel matrices.

Relations of bounded degree

Let $\mathcal{A} = \langle A, \bar{R} \rangle$ be a τ -structure.

We define a symmetric relation E_A on $\mathcal{A}$, and call $\langle A, E_A \rangle$ the **Gaifman-graph of $\mathcal{A}$** .

- Let $a, b \in A$. $(a, b) \in E_A$ iff there exists a relation $R \in \bar{R}$ and some $\bar{a} \in R$ such that both a and b appear in $\bar{a}$ (possibly with other members of A as well).
- For any element $a \in A$, the **degree** of a is the number of elements $b \neq a$ for which $(a, b) \in E_A$.
- We say that $\mathcal{A}$ is of **bounded degree d** if every $a \in A$ has degree at most d .
- We say that $\mathcal{A}$ is **connected** if its Gaifman-graph is connected.
- For a class of structures $\mathcal{P}$ we say it is of bounded degree d (resp. connected) iff all its structures are of bounded degree d (resp. connected).

The Specker-Blatter Theorem revisited, I

Theorem(E. Fischer and JAM, 2002)

Let $\mathcal{P}$ be a property of τ -structures, which is CMSOL-definable.
Let $d_{\mathcal{P}}(n)$ be its density function.

- If $\mathcal{P}$ is of bounded degree d ,
the function $d_{\mathcal{P}}(n)$ satisfies a modular recurrence relation for every m .
- Furthermore, if additionally all the models in $\mathcal{P}$ are connected,
the function $d_{\mathcal{P}}$ satisfies the trivial recurrence relations for every m .

We have **no restrictions on τ**
besides not allowing function symbols,

Hankel matrices for graph properties, I

- A k -graph of order n is a graph $G = (V(G), E(G), a_1, \dots, a_k)$ with $V(G) = [n]$ together with k distinct vertices $a_1, \dots, a_k$.
- Let $\square$ be a binary operation on k -graphs.
 $\square$ could be
 - (i) $\sqcup$, the disjoint union of graphs (without labels),
 - (ii) $\sqcup_k$, the k union of k -graphs formed by taking the disjoint union and then identifying corresponding distinct vertices.
 - (iii) or the substitution $S(G, H)$ of a 1-graph G for the distinct label a of the 1-graph H .

Hankel matrices for graph properties, II

- Let $G_i, i \in \mathbb{N}$ be an enumeration of all finite k -graphs.
- Let P be a graph property.
- We define the infinite $(0, 1)$ -matrix $\mathfrak{H}(P, \square)$ with

$$\mathfrak{H}(P, \square)_{i,j} = \begin{cases} 1 & G_i \square G_j \in P \\ 0 & G_i \square G_j \notin P \end{cases}$$

- We denote by $rk(P, \square)$ the rank of $\mathfrak{H}(P, \square)$ over the field $\mathbb{Z}^2$.

The matrix $\mathfrak{H}(P, \square)$ is called the **Hankel matrix** of P and $\square$ for k -graphs.

Unknowingly, C. Blatter and E. Specker use
a special case of Hankel matrices.

Hankel matrices (over a field $\mathcal{F}$)

Let $f : \mathcal{F} \rightarrow \mathcal{F}$ be a function over a field $\mathcal{F}$.

A **finite** or **infinite** matrix $H(f) = h_{i,j}$ is a **Hankel matrix for f** if $h_{i,j} = f(i+j)$.

Hankel matrices have **many** applications in:
numeric analysis, **probability theory** and **combinatorics**.

- Padé approximations
- Orthogonal polynomials
- Probability theory (theory of moments)
- Coding theory (BCH codes, Berlekamp-Massey algorithm)
- Combinatorial enumerations
(Lattice paths, Young tableaux, matching theory)

Hankel matrices over words

Let Σ be a finite alphabet and $\mathcal{F}$ be a field and let $f : \Sigma^* \rightarrow \mathcal{F}$ be a function on words.

A **finite** or **infinite** matrix $H(f) = h_{u,v}$ indexed over the words $u, v \in \Sigma^*$ is a **Hankel matrix for f** if $h_{u,v} = f(u \circ v)$. Here $\circ$ denotes concatenation.

Hankel matrices over words have applications in

- Formal language theory and stochastic automata,
J. Carlyle and A. Paz 1971
- Learning theory (exact learning of queries).
A. Beimel, F. Bergadano, N. Bshouty, E. Kushilevitz, S. Varricchio 1998
J. Oncina 2008
- Definability of picture languages.
O. Matz 1998, and D. Giammarresi and A. Restivo 2008

Hankel matrices for graphs

If we want to define Hankel matrices for (labeled) graphs,

what plays the role of concatenation?

- Disjoint union
Used by Freedman, Lovász and Schrijver, 2007, for characterizing multiplicative graph parameters over the real numbers
- k -unions (connections, connection matrices)
Used by Freedman, Lovász, Schrijver and Szegedy, 2007ff, for characterizing various forms and partition functions.
- Joins, cartesian products, generalized sum-like operations
used by Godlin, Kotek and JAM (2008) to prove non-definability.

The $\square$ -rank of a property P

Given

- a binary operation $\square$ between structures $\mathfrak{A}(A, \bar{R}, \bar{F}, \bar{c})$ with relations $\bar{R}$, functions $\bar{F}$ and constants $\bar{c}$, and no restrictions on their number of arguments, and
- a property P of such structures

we say that P has $\square$ -rank $r(P, \square) \in \mathbb{N} \cup \{\infty\}$ if $r(P, \square) = rk(P, \square)$ is the rank of the Hankel matrix $\mathfrak{H}(P, \square)$ over the field $\mathbb{Z}^2$.

$\mathcal{L}$ -smooth operations.

Let $\mathcal{L}$ be a logic with a well defined quantifier rank.

Here $\mathcal{L}$ is one of **FOL**, **MSOL**, **CMSOL**.

We say that two graphs G, H are $(\mathcal{L}, q)$ -equivalent, and write $G \sim_{\mathcal{L}}^q H$, if G and H satisfy the same $\mathcal{L}$ -sentences of quantifier rank q .

We say that $\square$ is $\mathcal{L}$ -smooth, if whenever we have

$$G_i \sim_{\mathcal{L}}^q H_i, i = 0, 1$$

then

$$G_0 \square G_1 \sim_{\mathcal{L}}^q H_0 \square H_1$$

This definition can be adapted to k -graphs and also to structures with k -ary operations for $k \geq 1$.

Proving that an operation $\square$ is $\mathcal{L}$ -smooth

Proving that an operation $\square$ is $\mathcal{L}$ -smooth may be difficult.

For **FOL**, **MSOL** and **CMSOL** this can be achieved using

Ehrenfeucht-Fraïssé games also know as **pebble games**.

The case for **MSOL** was first proved by Hans Läuchli in 1966.

Another way of establishing smoothness is via generalizations of the **Feferman-Vaught theorem**.

Makowsky, Johann A. "Algorithmic uses of the FefermanVaught theorem."
Annals of Pure and Applied Logic 126.1-3 (2004): 159-213.

In this talk we need two cases:

$\square = \sqcup$, the disjoint union, and $\square = \textit{Subst}$, the substitution of pointed graphs.

Theorem:(EF and JAM) Both $\sqcup$ and $\textit{Subst}$ are **CMSOL**-smooth.

The Finite Rank Theorem

Theorem (Godlin, Kotek, Makowsky 2008):

Let f be a **numeric graph parameter** or **graph polynomial** for τ -structures **definable in $\mathcal{L}$** and taking values in an integral domain $\mathcal{R}$.

Let $\square$ be an $\mathcal{L}$ -smooth operation.

Then the Hankel matrix $\mathfrak{H}(f, \square)$ has **finite rank** over $\mathcal{R}$.

The **Proof** uses a Feferman-Vaught-type theorem for graph polynomials, due to **B. Courcelle, J.A.M. and U. Rotics, 2000**.

C. Blatter and E. Specker prove it for *Subst* and **MSOL**.

L. Lovasz proves it for partition functions.

It has many applications in graph theory.

Bounded $\sqcup$ -rank and bounded degree

Theorem(EF and JAM, 2002)

Let P be a property of structures over some fixed vocabulary.

If P has finite $\sqcup$ -rank and

all its members of are bounded degree d . Then

- $d_P^m(n)$ is ultimately periodic for every m .
- Furthermore, if additionally all the models in P are connected, $d_P^m(n) = 0$ for all m and n .

Example $d_{EQ_2}(n)$ shows that bounded degree cannot be dropped, even for $\sqcup$ -rank 2 (and connected structures).

For structures of unbounded degree one needs a stronger assumption, the finiteness of the *Subst*-rank.

$\sqcup$ -rank of $\mathcal{P}$ and Gessel-classes

- A class of structures $\mathcal{P}$ is a **Gessel class** if for every $\mathfrak{A}$ and $\mathfrak{B}$, $\mathfrak{A} \sqcup \mathfrak{B} \in \mathcal{P}$ iff both $\mathfrak{A} \in \mathcal{P}$ and $\mathfrak{B} \in \mathcal{P}$.
- The class of forests is a Gessel class.
- If $\mathcal{P}$ is hereditary and closed under disjoint unions, it is a Gessel class.
- Every Gessel class has $\sqcup$ -rank at most 2.
- If $\mathcal{P}$ is a class of connected graphs, $\mathcal{P}$ has $\sqcup$ -rank at most 2, but is not a Gessel class.
- If $\mathcal{P}_1$ and $\mathcal{P}_2$ have finite $\sqcup$ -rank, so do $\mathcal{P}_1 \cup \mathcal{P}_2$, $\mathcal{P}_1 \cap \mathcal{P}_2$, and the complement $\bar{\mathcal{P}}_1$.

Many classes of finite $\sqcup$ -rank and *Subst*-rank.

There are only countably many classes of structures definable by **CMSOL**-formulas.

How many classes are there of finite $\sqcup$ -rank?

Proposition:(After an idea of Specker, 2002)

- (i) There are continuum many Gessel classes.
- (ii) There are continuum many properties of structures with $\sqcup$ -rank 2.
- (iii) There are continuum many properties of structures with finite *Subst*-rank.

Relation between $\sqcup$ -rank and *Subst*-rank

We denote by $G \bowtie H$ the join of G and H .

Proposition: Let P be a graph property. We have

- (i) $rk(P, \sqcup) \leq rk(P, Subst)$ and $rk(P, \bowtie) \leq rk(P, Subst)$.
- (ii) The class HAM of Hamiltonian graphs is connected,
hence has $\sqcup$ -rank 2,
However, HAM has infinite $\bowtie$ -rank, hence also infinite *Subst*-rank.

Problem: What can we say about $d_{HAM}^m(n)$?

The widely overlooked paper: its merits

- It initiates a systematic study of congruence relations for combinatorial counting functions.
Before, and still now, mostly special cases are studied.
- It anticipates the use of the rank of Hankel matrices in combinatorics
In the recent work of L. Lovász and his many collaborators on **graph limits** and **partition functions** this plays a central role.
- It formulates for the first time a **meta-theorem** connecting definability in **MSOL** with combinatorics.
The Büchi-Elgot-Trakhtenbrot Theorem characterizes **regular languages** using MSOL.
Bruno Courcelle and his collaborators use MSOL-definability to show that NP-hard problems are in P for graphs properties of **bounded tree- and clique-width**.

The widely overlooked paper: Our improvements

- Eldar Fischer gave a [counter example](#) using a [quaternary relation](#).
- Analyzing the counter example we noticed the role of [bounded degree structures](#) and the role of the [\$\sqcup\$ -rank](#).
- Proving an earlier version of the Finite Rank Theorem for **CMSOL** we extended the Specker-Blatter Theorem to include [Eulerian graphs](#) (and many other cases).

Thank you for your attention

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